Large Scale Density-friendly Graph Decomposition via Convex Programming

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The density-friendly decomposition (Tatti and Gionis WWW2015) is interesting as it merges two classic graph mining concepts: (i) the k-core decomposition and (ii) the notion of dense subgraph.

Theoretical analysis via convex programing

Given an edge-weighted (hyper) graph G = (V, E, w), we consider the following quadratic convex programing.

$$CP(G)$$

min $\sum_{u \in V} r_u^2$
s.t. $\forall u \in V, r_u = \sum_{e:u \in e} \alpha_u^e$
 $\forall e \in E, \sum_{u \in e} \alpha_u^e = w_e$
 $\forall u \in e \in E, \alpha_u^e \ge 0$



The Frank-Wolfe algorithm is a **projection free gradient-descent** method which has **convergence guarantees** for convex problems.

Contributions

Density-friendly is now as intuitive and almost as fast to compute as PageRank.

- We gave a new and **intuitive definition** of the density-friendly decomposition.
- We made an interesting link between the density-friendly decomposition and convex programing.
- We scaled up the computation of the density-friendly decomposition using the Frank-Wolfe algorithm.

Full paper at WWW2017. Code in C: https://github.com/maxdan94

Definition

Collection of non-overlapping sets of nodes $\{B_i\}$, such that

- B_1 maximizes $\frac{e_1}{n_1}$ and has maximum size (B_1 is the densest subgraph),
- B_2 maximizes $\frac{e_2+e_{12}}{n_2}$ and has maximum size,
- B_3 maximizes $\frac{e_3+e_{13}+e_{23}}{n_2}$ and has maximum size, ...



Theorem: our definition is equivalent to the one of Tatti and Gionis WWW2015.



Applying Frank-Wolfe on our quadratic convex programing leads to an algorithm very similar to our very simple algorithm.

Theorem (correctness): The level sets of an optimal solution r^G to the convex programing give the density-friendly decomposition (r^G gives the densities). **Theorem** (running time): After $t > \frac{4\Delta|E|}{\epsilon^2}$ iterations we have $||r^{(t)} - r^G||_2 \leq \epsilon$.

Experiments

Table 1: Our set of large graphs.

[1	Tatti and Cionis M/M/M/2015				
networks	n	m				
LiveJournal	4,036,538	34,681,189		Networks	exact	
Wikipedia	2,080,370	42,336,692		LiveJournal	2m45s	
Orkut	3,072,627	117,185,083		Wikipedia	2m14s	
Twitter	52,579,683	$1,\!614,\!106,\!500$		Orkut	13m08s	
Friendster	124,836,180	$1,\!806,\!067,\!135$		Twitter	4h57m28s	
gsh-2015	988,490,691	25,690,705,119		Friendster	5h48m27s	

Table 2: Running time comparison of our exact algorithm to the maxflow algorithm of WWW2015.

TG15

12m02s

7m07s

1h02m23s

gsh-2015	988,490,691	25,690,705,119		Friendster	5h48m27s	_					
We use a machine with 64G of RAM for all networks except gsh-2015 for which											
we use a machine with 512G of RAM.											

- Convergence in practice much faster than the worst case one
- 10^{-3} approximation of the densest subgraph within 300 iterations

A very simple, yet powerful, algorithm



Theorem: rank the nodes in decreasing order of density score:

- for t large enough, the nodes in the densest subgraph are ranked first;
- for t large enough, we have a density-friendly ordering.

- 10^{-2} approximation of the full decomposition within 1000 iterations on all networks except gsh-2015 (almost 10^{-1})
- the densest subgraph in gsh-2015 (25G edges) within 10 hours of computation



This research was partially supported by the Hong Kong RGC under the grants 17202715 and 17217716.

This research was partially supported by French National Agency (ANR) under project FIELDS (ANR-15-CE23-0006) and by a Google Faculty Award.